# Memory and Thread Placement Optimization Developer's Guide

Beta



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### **Preface**

The *Memory and Thread Placement Optimization Developer's Guide* provides information on locality groups and the technologies that are available to optimize the use of computing resources in the Oracle Solaris operating system.

### Who Should Use This Book

This book is intended for use by system administrators, performance engineers, systems programmers, and support engineers, and developers who are writing applications in an environment with multiple CPUs and a non-uniform memory architecture. The programming interfaces and tools that are described in this book give the developer control over the system's behavior and resource allocation.

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Third-party URLs are referenced in this document and provide additional, related information.

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### **Typographic Conventions**

The following table describes the typographic conventions that are used in this book.

TABLE P-1 Typographic Conventions

Typeface	Meaning	Example
AaBbCc123	The names of commands, files, and directories, and onscreen computer output	Edit your . login file.
		Use ls -a to list all files.
		<pre>machine_name% you have mail.</pre>
AaBbCc123	What you type, contrasted with onscreen	machine_name% <b>su</b>
	computer output	Password:
aabbcc123	Placeholder: replace with a real name or value	The command to remove a file is rm <i>filename</i> .
AaBbCc123	Book titles, new terms, and terms to be	Read Chapter 6 in the <i>User's Guide</i> .
	emphasized	A <i>cache</i> is a copy that is stored locally.
		Do <i>not</i> save the file.
		<b>Note:</b> Some emphasized items appear bold online.

# **Shell Prompts in Command Examples**

The following table shows the default UNIX system prompt and superuser prompt for shells that are included in the Oracle Solaris OS. Note that the default system prompt that is displayed in command examples varies, depending on the Oracle Solaris release.

TABLE P-2 Shell Prompts

Shell	Prompt
Bash shell, Korn shell, and Bourne shell	\$
Bash shell, Korn shell, and Bourne shell for superuser	#
C shell	machine_name%
C shell for superuser	machine_name#

◆ ◆ ◆ CHAPTER 1

## Overview of Locality Groups

- "Locality Groups Overview" on page 9
- "MPO Observability Tools" on page 12

### **Locality Groups Overview**

Shared memory multiprocessor computers contain multiple CPUs. Each CPU can access all of the memory in the machine. In some shared memory multiprocessors, the memory architecture enables each CPU to access some areas of memory more quickly than other areas.

When a machine with such a memory architecture runs the Oracle Solaris software, providing information to the kernel about the shortest access times between a given CPU and a given area of memory can improve the system's performance. The locality group (lgroup) abstraction has been introduced to handle this information. The lgroup abstraction is part of the Memory Placement Optimization (MPO) feature.

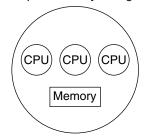
An Igroup is a set of CPU-like and memory-like devices in which each CPU in the set can access any memory in that set within a bounded latency interval. The value of the latency interval represents the least common latency between all the CPUs and all the memory in that Igroup. The latency bound that defines an Igroup does not restrict the maximum latency between members of that Igroup. The value of the latency bound is the shortest latency that is common to all possible CPU-memory pairs in the group.

Lgroups are hierarchical. The lgroup hierarchy is a Directed Acyclic Graph (DAG) and is similar to a tree, except that an lgroup might have more than one parent. The root lgroup contains all the resources in the system and can include child lgroups. Furthermore, the root lgroup can be characterized as having the highest latency value of all the lgroups in the system. All of its child lgroups will have lower latency values. The lgroups closer to the root have a higher latency while lgroups closer to leaves have lower latency.

A computer in which all the CPUs can access all the memory in the same amount of time can be represented with a single lgroup (see Figure 1-1). A computer in which some of the CPUs can access some areas of memory in a shorter time than other areas can be represented by using multiple lgroups (see Figure 1-2).

FIGURE 1-1 Single Locality Group Schematic

Machine with single latency is represented by one Igroup



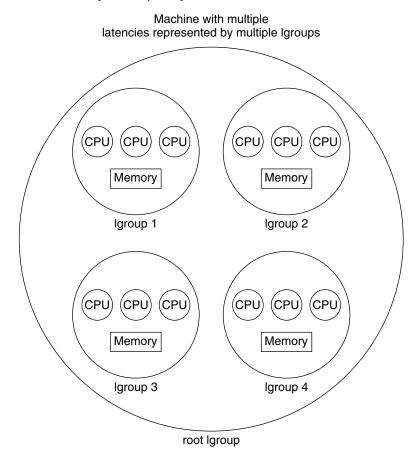


FIGURE 1–2 Multiple Locality Groups Schematic

The organization of the Igroup hierarchy simplifies the task of finding the nearest resources in the system. Each thread is assigned a home Igroup upon creation. The operating system attempts to allocate resources for the thread from the thread's home Igroup by default. For example, the Oracle Solaris kernel attempts to schedule a thread to run on the CPUs in the thread's home Igroup and allocate the thread's memory in the thread's home Igroup by default. If the desired resources are not available from the thread's home Igroup, the kernel can traverse the Igroup hierarchy to find the next nearest resources from parents of the home Igroup. If the desired resources are not available in the home Igroup's parents, the kernel continues to traverse the Igroup hierarchy to the successive ancestor Igroups of the home Igroup. The root Igroup is the ultimate ancestor of all other Igroups in a machine and contains all of the machine's resources.

The Memory Placement Optimization (MPO) tools enable developers to tune the performance of the MPO features in cases where the default MPO behaviors do not yield the desired performance.

The Igroup APIs export the Igroup abstraction for applications to use for observability and performance tuning. A new library, called liblgrp, contains the new APIs. Applications can use the APIs to perform the following tasks:

- Traverse the group hierarchy
- Discover the contents and characteristics of a given lgroup
- Affect the thread and memory placement on Igroups

### **MPO Observability Tools**

The MPO tools help developers to answer questions about system configuration and balance or placement. The tools also provide the basic information and mechanisms that developers need in order to determine whether MPO is successful and to diagnose problems related to MPO.

To determine the degree of success that MPO has in providing useful locality assignments and acceptable performance, it is important to know a given thread's affinities for lgroups, including its home lgroup, and where the thread's memory is allocated.

The MPO observability tools provide developers with the ability to determine the actions taken by the system. The MPO thread and memory placement tools enable developers to act on that information. Developers can also use the dtrace(1M) tool to gain further insights into the system's behavior.



# MPO Observability Tools

This chapter describes the tools that are available to use the MPO functionality that is available in the Oracle Solaris operating system.

This chapter discusses the following topics:

- "The pmadvise utility" on page 13 describes the tool that applies rules that define the memory use of a process.
- "Using the madv.so.1 Shared Object" on page 15 describes the madv.so.1 shared object and how to use it to configure virtual memory advice.
- "The plgrp tool" on page 17 describes the tool that can display and set a thread's affinity for a locality group.
- "The lgrpinfo Tool" on page 19 prints information about the lgroup hierarchy, contents, and characteristics.
- "The Solaris::lgrp Module" on page 21 describes a Perl interface to the locality group API that is described in Chapter 3, "Locality Group APIs."

### The pmadvise utility

The pmadvise utility applies rules to a process that define how that process uses memory. The pmadvise utility applies the rules, called *advice*, to the process with the madvise(3C) tool. This tool can apply advice to a specific subrange of locations in memory at a specific time. By contrast, the madv.so.1(1) tool applies the advice throughout the execution of the target program to all segments of a specified type.

The pmadvise utility has the following options:

- -f This option takes control of the target process. This option overrides the control of any other process. See the proc(1) manual page.
- This option specifies the advice to apply to the target process. Specify the advice in this format:

private=advice
shared=advice
heap=advice
stack=advice
address:length=advice

The value of the *advice* term can be one of the following:

normal random sequential willneed dontneed free access\_lwp access\_many access default

You can specify an address and length to specify the subrange where the advice applies. Specify the address in hexadecimal notation and the length in bytes.

If you do not specify the length and the starting address refers to the start of a segment, the pmadvise utility applies the advice to that segment. You can qualify the length by adding the letters K, M, G, T, P, or E to specify kilobytes, megabytes, gigabytes, terabytes, or exabytes, respectively.

-v This option prints verbose output in the style of the pmap(1) tool that shows the value and locations of the advice rules currently in force.

The pmadvise tool attempts to process all legal options. When the pmadvise tool attempts to process an option that specifies an illegal address range, the tool prints an error message and skips that option. When the pmadvise tool finds a syntax error, it quits without processing any options and prints a usage message.

When the advice for a specific region conflicts with the advice for a more general region, the advice for the more specific region takes precedence. Advice that specifies a particular address range has precedence over advice for the heap and stack regions, and advice for the heap and stack regions has precedence over advice for private and shared memory.

The advice rules in each of the following groups are mutually exclusive from other advice rules within the same group:

```
MADV_NORMAL, MADV_RANDOM, MADV_SEQUENTIAL
MADV_WILLNEED, MADV_DONTNEED, MADV_FREE
MADV_ACCESS_DEFAULT, MADV_ACCESS_LWP, MADV_ACCESS_MANY
```

### Using the madv.so.1 Shared Object

The madv.so.1 shared object enables the selective configuration of virtual memory advice for launched processes and their descendants. To use the shared object, the following string must be present in the environment:

LD\_PRELOAD=\$LD\_PRELOAD:madv.so.1

The madv.so.1 shared object applies memory advice as specified by the value of the MADV environment variable. The MADV environment variable specifies the virtual memory advice to use for all heap, shared memory, and mmap regions in the process address space. This advice is applied to all created processes. The following values of the MADV environment variable affect resource allocation among lgroups:

access\_default This value resets the kernel's expected access pattern to the default.

This value advises the kernel that the next LWP to touch an address range is the LWP that will access that range the most. The kernel allocates the memory and other resources for this range and the LWP accordingly.

This value advises the kernel that many processes or LWPs will access memory randomly across the system. The kernel allocates the memory

and other resources accordingly.

The value of the MADVCFGFILE environment variable is the name of a text file that contains one

The value of *exec-name* is the name of an application or executable. The value of *exec-name* can be a full pathname, a base name, or a pattern string.

or more memory advice configuration entries in the form *exec-name:advice-opts*.

The value of *advice-opts* is of the form *region=advice*. The values of *advice* are the same as the values for the MADV environment variable. Replace *region* with any of the following legal values:

madv	Advice applies to all heap, shared memory, and $mmap(2)$ regions in the process address space.
heap	The heap is defined to be the brk(2) area. Advice applies to the existing heap and to any additional heap memory allocated in the future.
shm	Advice applies to shared memory segments. See shmat(2) for more information on shared memory operations.
ism	Advice applies to shared memory segments that are using the SHM_SHARE_MMU flag. The ism option takes precedence over shm.
dsm	Advice applies to shared memory segments that are using the SHM_PAGEABLE flag. The dsm option takes precedence over shm.

mapshared Advice applies to mappings established by the mmap() system call using the

MAP\_SHARED flag.

mapprivate Advice applies to mappings established by the mmap() system call using the

MAP PRIVATE flag.

mapanon Advice applies to mappings established by the mmap() system call using the

MAP ANON flag. The mapanon option takes precedence when multiple options

apply.

The value of the MADVERRFILE environment variable is the name of the path where error messages are logged. In the absence of a MADVERRFILE location, the madv. so. 1 shared object logs errors by using syslog(3C) with a LOG\_ERR as the severity level and LOG\_USER as the facility descriptor.

Memory advice is inherited. A child process has the same advice as its parent. The advice is set back to the system default advice after a call to exec(2) unless a different level of advice is configured using the madv.so.1 shared object. Advice is only applied to mmap() regions explicitly created by the user program. Regions established by the run-time linker or by system libraries that make direct system calls are not affected.

### madv.so.1 Usage Examples

The following examples illustrate specific aspects of the madv. so. 1 shared object.

EXAMPLE 2-1 Setting Advice for a Set of Applications

This configuration applies advice to all ISM segments for applications with exec names that begin with foo.

**EXAMPLE 2-2** Excluding a Set of Applications From Advice

This configuration sets advice for all applications with the exception of ls.

```
$ LD_PRELOAD=$LD_PRELOAD:madv.so.1
$ MADV=access_many
$ MADVCFGFILE=madvcfg
$ export LD_PRELOAD MADV MADVCFGFILE
$ cat $MADVCFGFILE
$ ls:
```

#### EXAMPLE 2-3 Pattern Matching in a Configuration File

Because the configuration specified in MADVCFGFILE takes precedence over the value set in MADV, specifying \* as the *exec-name* of the last configuration entry is equivalent to setting MADV. This example is equivalent to the previous example.

#### EXAMPLE 2-4 Advice for Multiple Regions

This configuration applies one type of advice for mmap() regions and different advice for heap and shared memory regions for applications whose exec() names begin with foo.

### The plgrp tool

The plgrp utility can display or set the home lgroup and lgroup affinities for one or more processes, threads, or lightweight processes (LWPs). The system assigns a home lgroup to each thread on creation. When the system allocates a CPU or memory resource to a thread, it searches the lgroup hierarchy from the thread's home lgroup for the nearest available resources to the thread's home.

The system chooses a home lgroup for each thread. The thread's affinity for its home lgroup is initially set to none, or no affinity. When a thread sets an affinity for an lgroup in its processor set that is higher than the thread's affinity for its home lgroup, the system moves the thread to that lgroup. The system does not move threads that are bound to a CPU. The system rehomes a thread to the lgroup in its processor set that has the highest affinity when the thread's affinity for its home lgroup is removed (set to none).

For a full description of the different levels of Igroup affinity and their semantics, see the Igrp affinity set(3LGRP) manual page.

The plgrp tool supports the following options:

```
-a lgroup list
```

This option displays the affinities of the processes or threads that you specify for the lgroups in the list. - Algroup list/none | weak | strong[, . . . ] This option sets the affinity of the processes or threads that you specify for the lgroups in the list. You can use a comma separated list of *lgroup/affinity* assignments to set several affinities at once. - F This option takes control of the target process. This option overrides the control of any other process. See the proc(1) manual page. -h This option returns the home Igroup of the processes or threads that you specify. This is the default behavior of the plgrp tool when you do not specify any options. -H lgroup list This option sets the home Igroup of the processes or threads that you specify. This option sets a strong affinity for the listed lgroup. If you specify more than one Igroup, the plgrp utility will attempt to home the threads to the Igroups in a

### **Specifying Lgroups**

The value of the *lgroup list* variable is a comma separated list of one or more of the following attributes:

round robin fashion.

- lgroup ID
- Range of Igroup IDs, specified as start Igroup ID-end Igroup ID
- all
- root
- leaves

The all keyword represents all of the Igroup IDs in the system. The root keyword represents the ID of the root Igroup. The leaves keyword represents the IDs of all of the leaf Igroups. A leaf Igroup is an Igroup that does not have any children.

### **Specifying Process and Thread Arguments**

The plgrp utility takes one or more space-separated processes or threads as arguments. You can specify processes and threads in a the same syntax that the proc(1) tools use. You can specify a process ID as an integer, with the syntax *pid* or /proc/*pid*. You can use shell expansions with the /proc/*pid* syntax. When you give a process ID alone, the arguments to the plgrp utility include all of the threads of that process.

You can specify a thread explicitly by specifying the process ID and thread ID with the syntax *pid/lwpid*. You can specify multiple threads of a process by defining ranges with can be selected at once by using the - character to define a range, or with a comma-separated list. To specify threads 1, 2, 7, 8, and 9 of a process whose process ID is *pid*, use the syntax *pid/*1, 2, 7-9.

### The Igrpinfo Tool

The lgrpinfo tool prints information about the lgroup hierarchy, contents, and characteristics. The lgrpinfo tool is a Perl script that requires the Solaris::Lgrp module. This tool uses the liblgrp(3LIB) API to get the information from the system and displays it in the human-readable form.

The lgrpinfo tool prints general information about all of the lgroups in the system when you call it without any arguments. When you pass lgroup IDs to the lgrpinfo tool at the command line, the tool returns information about the lgroups that you specify. You can specify lgroups with their lgroup IDs or with one of the following keywords:

This keyword specifies all Igroups and is the default behavior.

root This keyword specifies the root lgroup.

leaves This keyword specifies all of the leaf lgroups. A leaf lgroup is an lgroup that

has no children in the lgroup hierarchy.

intermediate This keyword specifies all of the intermediate Igroups. An intermediate

lgroup is an lgroup that has a parent and children.

When the lgrpinfo tool receives an invalid lgroup ID, the tool prints a message with the invalid ID and continues processing any other lgroups that are passed in the command line. When the lgrpinfo tool finds no valid lgroups in the arguments, it exits with a status of 2.

### Options for the **lgrpinfo** Tool

When you call the lgrpinfo tool without any arguments, the tool's behavior is equivalent to using the options -celmrt all. The valid options for the lgrpinfo tool are:

- This option prints the topology, CPU, memory, load and latency information for the specified Igroup IDs. This option combines the behaviors of the -tcemrlL options, unless you also specify the -T option. When you specify the -T option, the behavior of the -a option does not include the behavior of the -t option.
- -c This option prints the CPU information.

- -C This option replaces each lgroup in the list with its children. You cannot combine this option with the -P or -T options. When you do not specify any arguments, the tool applies this option to all lgroups.
- -e This option prints Igroup load averages for leaf Igroups.
- This option prints the OS view of the lgroup hierarchy. The tool's default behavior displays the caller's view of the lgroup hierarchy. The caller's view only includes the resources that the caller can use. See the lgrp\_init(3LGRP) manual page for more details on the OS and caller's view.
- -h This option prints the help message for the tool.
- -I This option prints only IDs that match the IDs you specify. You can combine this option with the -c, -G, -C, or -P options. When you specify the -c option, the tool prints the list of CPUs that are in all of the matching lgroups. When you do not specify the -c option, the tool displays the IDs for the matching lgroups. When you do not specify any arguments, the tool applies this option to all lgroups.
- This option prints information about Igroup latencies. The latency value given for each Igroup is defined by the operating system and is platform-specific. It can only be used for relative comparison of Igroups on the running system. It does not necessarily represent the actual latency between hardware devices and may not be applicable across platforms.
- -L This option prints the lgroup latency table. This table shows the relative latency from each lgroup to each of the other lgroups.
- -m This option prints memory information. The tool reports memory sizes in the units that give a size result in the integer range from 0 to 1023. You can override this behavior by using the -u option. The tool will only display fractional results for values smaller than 10.
- -P This option replaces each lgroup in the list with its parent or parents. You cannot combine this option with the -C or -T options. When you do not specify any arguments, the tool applies this option to all lgroups.
- r This option prints information about lgroup resources. When you specify the -T option, the tool displays information about the resources of the intermediate lgroups only.
- -t This option prints information about Igroup topology.
- -T This option prints the Igroup topology of a system graphically, as a tree. You can only use this option with the -a, -c, -e, -G, -l, -L, -m, -r, and -u options. To restrict the output to intermediate Igroups, use the -r option. Omit the -t option when you combine the -T option with the -a option. This option does not print information for the root Igroup unless it is the only Igroup.

-uunits

This option specifies memory units. The value of the *units* argument can be b, k, m, g, t, p, or e for bytes, kilobytes, megabytes, gigabytes, terabytes, petabytes, or exabytes, respectively.

### The Solaris::lgrp Module

This Perl module provides a Perl interface to the Igroup APIs that are in liblgrp. This interface provides a way to traverse the Igroup hierarchy, discover its contents and characteristics, and set a thread's affinity for an Igroup. The module gives access to various constants and functions defined in the Igrp\_user.h header file. The module provides the procedural interface and the object interface to the library.

The default behavior of this module does not export anything. You can use the following tags to selectively import the constants and functions that are defined in this module:

:LGRP\_CONSTANTS LGRP\_AFF\_NONE, LGRP\_AFF\_STRONG, LGRP\_AFF\_WEAK,

LGRP\_CONTENT\_DIRECT, LGRP\_CONTENT\_HIERARCHY, LGRP\_MEM\_SZ\_FREE,

LGRP\_MEM\_SZ\_INSTALLED, LGRP\_VER\_CURRENT, LGRP\_VER\_NONE, LGRP\_VIEW\_CALLER, LGRP\_VIEW\_OS, LGRP\_NONE, LGRP\_RSRC\_CPU, LGRP\_RSRC\_MEM, LGRP\_CONTENT\_ALL, LGRP\_LAT\_CPU\_TO\_MEM

:PROC CONSTANTS P PID, P LWPID, P MYID

:CONSTANTS :LGRP CONSTANTS,:PROC CONSTANTS

:FUNCTIONS lgrp affinity get(),lgrp affinity set(),lgrp children(),

lgrp cookie stale(),lgrp cpus(),lgrp fini(),lgrp home(),

lgrp init(), lgrp latency(), lgrp latency cookie(),

lgrp\_mem\_size(), lgrp\_nlgrps(), lgrp\_parents(), lgrp\_root(),
lgrp\_version(), lgrp\_view(), lgrp\_resources(), lgrp\_lgrps(),

lgrp leaves(), lgrp isleaf(), lgrp lgrps(), lgrp leaves().

:ALL() :CONSTANTS(),:FUNCTIONS()

The Perl module has the following methods:

- new()
- cookie()
- stale()
- view()
- root()
- children()
- parents()
- nlgrps()
- mem size()
- cpus()

- isleaf()
- resources()
- version()
- home()
- affinity\_get()
- affinity set()
- lgrps()
- leaves()
- latency()

You can export constants with the : CONSTANTS or : ALL tags. You can use any of the constants in the following list in Perl programs.

- LGRP NONE
- LGRP\_VER\_CURRENT
- LGRP VER NONE
- LGRP VIEW CALLER
- LGRP VIEW OS
- LGRP AFF NONE
- LGRP AFF STRONG
- LGRP AFF WEAK
- LGRP CONTENT DIRECT
- LGRP CONTENT HIERARCHY
- LGRP MEM SZ FREE
- LGRP MEM SZ INSTALLED
- LGRP RSRC CPU
- LGRP RSRC MEM
- LGRP CONTENT ALL
- LGRP\_LAT\_CPU TO MEM
- P PID
- P LWPID
- P MYID

When an underlying library function fails, the functions in this module return either undef or an empty list. The module can use the following error codes:

EINVAL The value supplied is not valid.

ENOMEM There was not enough system memory to complete an operation.

ESRCH The specified process or thread was not found.

The effective user of the calling process does not have the appropriate privileges,

and its real or effective user ID does not match the real or effective user ID of one of

the threads.

### Functions in the Solaris::lgrp Module

#### lgrp\_init([LGRP\_VIEW\_CALLER | LGRP\_VIEW\_OS])

This function initializes the Igroup interface and takes a snapshot of the Igroup hierarchy with the given view. Given the view, the <code>lgrp\_init()</code> function returns a cookie that represents this snapshot of the Igroup hierarchy. Use this cookie with the other routines in the Igroup interface that require the Igroup hierarchy. Call the <code>lgrp\_fini()</code> function with this cookie when the system no longer needs the hierarchy snapshot. Unlike the <code>lgrp\_init(3LGRP)</code> function, this function assumes a value of <code>LGRP\_VIEW\_OS</code> as the default if the system provides no view. This function returns a cookie upon successful completion. If the <code>lgrp\_init</code> function does not complete successfully, it returns a value of undef and sets <code>\$!</code> to indicate the error. See the man page for the <code>lgrp\_init(3LGRP)</code> function for more information.

#### lgrp fini(\$cookie)

This function takes a cookie, frees the snapshot of the lgroup hierarchy that the lgrp\_init() function created, and cleans up anything else that the lgrp\_init() function set up. After calling this function, do not use the cookie that the lgroup interface returns. This function returns 1 upon successful completion. If the lgrp\_fini function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_fini(3LGRP) function for more information.

#### lgrp\_view(\$cookie)

This function takes a cookie that represents the snapshot of the lgroup hierarchy and returns the snapshot's view of the lgroup hierarchy. If the given view is LGRP\_VIEW\_CALLER, the snapshot contains only the resources that are available to the caller. When the view is LGRP\_VIEW\_OS, the snapshot contains the resources that are available to the operating system. This function returns the view for the snapshot of the lgroup hierarchy that is represented by the given cookie upon successful completion. If the lgrp\_view function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_view(3LGRP) function for more information.

#### lgrp home(\$idtype, \$id)

This function returns the home lgroup for the given process or thread. To specify a process, give the \$idtype argument the value P\_PID and give the \$id argument the value of the process id. To specify a thread, give the \$idtype argument the value P\_LWPID and give the \$id argument the value of the thread's LWP id. To specify the current process or thread, give the \$id argument the value P\_MYID. This function returns the id of the home lgroup of the specified process or thread upon successful completion. If the lgrp\_home function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_home(3LGRP) function for more information.

#### lgrp cookie stale(\$cookie)

This function returns the staleness status of the specified cookie upon successful completion. If the lgrp cookie stale function does not complete successfully, it returns a value of

undef and sets \$! to indicate the error. This function fails and returns EINVAL if the cookie is not valid. See the man page for the lgrp\_cookie\_stale(3LGRP) function for more information.

#### lgrp\_cpus(\$cookie, \$lgrp, \$context)

This function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the list of CPUs that are in the lgroup that is specified by the \$lgrp argument. Give the \$context argument the value LGRP\_CONTENT\_HIERARCHY to make the lgrp\_cpus() function return the list of all the CPUs that are in the specified lgroup, including child lgroups. Give the \$context() argument the value LGRP\_CONTENT\_DIRECT to make the lgrp\_cpus() function return the list of CPUs that are directly contained in the specified lgroup. This function returns the number of CPUs that are in the specified lgroup when you call it in a scalar context. If the lgrp\_cpus function does not complete successfully when you call it in a scalar context, it returns a value of undef and sets \$! to indicate the error. If the lgrp\_cpus function does not complete successfully when you call it in a list context, it returns the empty list and sets \$! to indicate the error. See the man page for the lgrp\_cpus(3LGRP) function for more information.

#### lgrp children(\$cookie, \$lgrp)

This function takes a cookie that represents a snapshot of the Igroup hierarchy and returns the list of Igroups that are children of the specified Igroup. When called in scalar context, the Igrp\_children() function returns the number of children Igroups for the specified Igroup when you call it in a scalar context. If the Igrp\_children function does not complete successfully when you call it in a scalar context, it returns a value of undef and sets \$! to indicate the error. If the Igrp\_children function does not complete successfully when you call it in a list context, it returns the empty list and sets \$! to indicate the error. See the man page for the Igrp\_children(3LGRP) function for more information.

#### lgrp parents(\$cookie, \$lgrp)

This function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the list of parent groups of the specified lgroup. When called in scalar context, the lgrp\_parents() function returns the number of parent lgroups for the specified lgroup when you call it in a scalar context. If the lgrp\_parents function does not complete successfully when you call it in a scalar context, it returns a value of undef and sets \$! to indicate the error. If the lgrp\_parents function does not complete successfully when you call it in a list context, it returns the empty list and sets \$! to indicate the error. See the man page for the lgrp\_parents(3LGRP) function for more information.

#### lgrp nlgrps(\$cookie)

This function takes a cookie that represents a snapshot of the lgroup hierarchy. It returns the number of lgroups in the hierarchy. This number is always at least one. If the lgrp\_nlgrps function does not complete successfully, it returns a value of undef and sets the value of \$! to EINVAL to indicate that the cookie is invalid. See the man page for the lgrp\_nlgrps(3LGRP) function for more information.

#### lgrp root(\$cookie)

This function returns the ID of the root lgroup. If the lgrp\_root function does not complete successfully, it returns a value of undef and sets the value of \$! to EINVAL to indicate that the cookie is invalid. See the man page for the lgrp\_root(3LGRP) function for more information.

#### lgrp\_mem\_size(\$cookie, \$lgrp, \$type, \$content)

This function takes a cookie that represents a snapshot of the lgroup hierarchy. The function returns the memory size of the given lgroup in bytes. Set the value of the \$type argument to LGRP\_MEM\_SZ\_FREE to have the lgrp\_mem\_size() function return the amount of free memory. Set the value of the \$type argument to LGRP\_MEM\_SZ\_INSTALLED to have the lgrp\_mem\_size() function return the amount of installed memory. Set the value of the \$content argument to LGRP\_CONTENT\_HIERARCHY to have the lgrp\_mem\_size() function return results for the specified lgroup and each of its child lgroups. Set the value of the \$content argument to LGRP\_CONTENT\_DIRECT to have the lgrp\_mem\_size() function return results for the specified lgroup only. This function returns the memory size in bytes upon successful completion, the size in bytes is returned. If the lgrp\_mem\_size function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_mem\_size(3LGRP) function for more information.

#### lgrp version([VERSION])

This function takes an interface version number as the value of the VERSION argument and returns an Igroup interface version. To discover the current Igroup interface version, pass the value of LGRP\_VER\_CURRENT or LGRP\_VER\_NONE in the VERSION argument. The Igrp\_version() function returns the requested version if the system supports that version. The Igrp\_version() function returns LGRP\_VER\_NONE if the system does not supports the request version. The Igrp\_version() function returns the current version of the library when you call the function with LGRP\_VER\_NONE as the value of the VERSION argument. This code fragment tests whether the version of the interface used by the caller is supported:

See the man page for the lgrp\_version(3LGRP) function for more information.

#### lgrp affinity set(\$idtype, \$id, \$lgrp, \$affinity)

This function sets the affinity that the LWPs you specify with the \$idtype and \$id arguments have for the given Igroup. You can set the Igroup affinity to LGRP\_AFF\_STRONG, LGRP\_AFF\_WEAK, or LGRP\_AFF\_NONE. When the value of the \$idtype argument is P\_PID, this function sets the affinity for all the LWPs of the process with the process id specified in the \$id argument. The lgrp\_affinity\_set() function sets the affinity for the LWP of the current process with LWP id \$id when the value of the \$idtype argument is P\_LWPID. You can specify the current LWP or process by assigning the \$id argument a value of P\_MYID. This function returns 1 upon successful completion. If the lgrp\_affinity\_set function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_affinity\_set(3LGRP) function for more information.

#### lgrp affinity get(\$idtype, \$id, \$lgrp)

This function retrieves the affinity that the LWPs you specify with the \$idtype and \$id arguments have for the given Igroup. When the value of the \$idtype argument is P\_PID, this function retrieves the affinity for one of the LWPs in the process. The lgrp\_affinity\_get() function retrieves the affinity for the LWP of the current process with LWP id \$id when the value of the \$idtype argument is P\_LWPID. You can specify the current LWP or process by assigning the \$id argument a value of P\_MYID. This function returns 1 upon successful completion. If the lgrp\_affinity\_get function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_affinity\_get(3LGRP) function for more information.

lgrp\_latency\_cookie(\$cookie, \$from, \$to, [\$between=LGRP\_LAT\_CPU\_TO\_MEM])()
This function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the latency value between a hardware resource in the \$from lgroup to a hardware resource in the \$to lgroup. This function returns the latency value within a given lgroup when the values of the \$from and \$to arguments are identical. Set the value of the optional \$between argument to LGRP\_LAT\_CPU\_TO\_MEM to specify the hardware resources to measure the latency between. LGRP\_LAT\_CPU\_TO\_MEM represents the latency from CPU to memory and is the only valid value for this argument in this release. This function returns 1 upon successful completion. If the lgrp\_latency\_cookie function does not complete successfully, it returns a value of undef and sets \$! to indicate the error. See the man page for the lgrp\_latency\_cookie(3LGRP) function for more information.

#### lgrp\_latency(\$from, \$to)()

The function is similar to the <code>lgrp\_latency\_cookie()</code> function, but returns the latency between the given lgroups at the given instant in time. Because the system dynamically reallocates and frees lgroups, this function's results are not always consistent across calls. This function is deprecated. Use the <code>lgrp\_latency\_cookie()</code> function instead. See the man page for the <code>lgrp\_latency(3LGRP)</code> function for more information.

#### lgrp resources(\$cookie, \$lgrp, \$type)()

This function is only available for version 2 of the API. When you call this function with version 1 of the API, the lgrp\_resources() function returns undef or the empty list and sets the value of \$! to EINVAL. This function returns the list of lgroups that directly contain the specified type of resources. The resources are represented by a set of lgroups in which each lgroup directly contains CPU and/or memory resources. To specify CPU resources, set the value of the \$type argument to LGRP\_RSRC\_CPU. To specify memory resources, set the value of the \$type argument to LGRP\_RSRC\_MEM. If the lgrp\_resources function does not complete successfully, it returns a value of undef or the empty list and sets \$! to indicate the error. See the man page for the lgrp\_resources(3LGRP) function for more information.

#### lgrp\_lgrps(\$cookie, [\$lgrp])()

This function returns the list of all of the lgroups in a hierarchy, starting from the lgroup specified in the \$lgrp argument. This function uses the value returned by the lgrp\_root(\$cookie) function when the \$lgrp argument has no value. The lgrp\_lgrps() function returns the empty list on failure. This function returns the total number of lgroups in the system when you call it in a scalar context.

#### lgrp leaves(\$cookie, [\$lgrp])()

This function returns the list of all leaf lgroups in a hierarchy that starts from the lgroup specified in the \$lgrp argument. This function uses the value returned by the lgrp\_root(\$cookie) function when the \$lgrp argument has no value. The lgrp\_leaves() function returns undef or the empty list on failure. This function returns the total number of leaf lgroups in the system when you call it in a scalar context.

#### lgrp isleaf(\$cookie, \$lgrp)()

This function returns True if the lgroup specified by the value of the \$lgrp argument is a leaf lgroup. Leaf lgroups have no children. The lgrp\_isleaf() function returns False if the specified lgroup is not a leaf lgroup.

### Object Methods in the Solaris::lgrp Module

#### new([\$view])

This method creates a new Sun::Solaris::Lgrp object. An optional argument is passed to the lgrp\_init() function. This method uses a value for the \$view argument of LGRP VIEW OS by default.

#### cookie()

This function returns a transparent cookie that is passed to functions that accept a cookie.

#### version([\$version])

This method returns the current version of the liblgrp(3LIB) library when you call it without an argument. This is a wrapper for the lgrp\_version() function with LGRP VER NONE as the default value of the \$version argument.

#### stale()

This method returns T if the Igroup information in the object is stale. This method returns F in all other cases. The stale method is a wrapper for the Igrp cookie stale() function.

#### view()

This method returns the snapshot's view of the lgroup hierarchy. The view() method is a wrapper for the lgrp\_view() function.

#### root()

This method returns the root lgroup. The root() method is a wrapper for the lgrp\_root() function.

#### children(\$lgrp)

This method returns the list of Igroups that are children of the specified Igroup. The children method is a wrapper for the lgrp children() function.

#### parents(\$lgrp)

This method returns the list of Igroups that are parents of the specified Igroup. The parents method is a wrapper for the Igrp\_parents() function.

#### nlgrps()

This method returns the number of Igroups in the hierarchy. The nlgrps() method is a wrapper for the lgrp nlgrps() function.

#### mem size(\$lgrp, \$type, \$content)

This method returns the memory size of the given lgroup in bytes. The mem\_size method is a wrapper for the lgrp mem size() function.

#### cpus(\$lgrp, \$context)

This method returns the list of CPUs that are in the lgroup specified by the \$lgrp argument. The cpus method is a wrapper for the lgrp\_cpus() function.

#### resources(\$lgrp, \$type)

This method returns the list of Igroups that directly contain resources of the specified type. The resources method is a wrapper for the Igrp resources () function.

#### home(\$idtype, \$id)

This method returns the home lgroup for the given process or thread. The home method is a wrapper for the lgrp\_home() function.

#### affinity get(\$idtype, \$id, \$lgrp)

This method returns the affinity that the LWP has to a given lgroup. The affinity\_get() method is a wrapper for the lgrp affinity get() function.

#### affinity\_set(\$idtype, \$id, \$lgrp, \$affinity)

This method sets the affinities that the LWPs specified by the \$idtype and \$id arguments have for the given lgroup. The affinity\_set() method is a wrapper for the lgrp\_affinity\_set() function.

#### lgrps([\$lgrp])

This method returns the list of all of the lgroups in a hierarchy starting from the lgroup specified by the value of the \$lgrp argument. The hierarchy starts from the root lgroup when you do not specify a value for the \$lgrp argument. The lgrps() method is a wrapper for the lgrp lgrps() function.

#### leaves([\$lgrp])

This method returns the list of all of the leaf lgroups in a hierarchy starting from the lgroup specified by the value of the \$lgrp argument. The hierarchy starts from the root lgroup when you do not specify a value for the \$lgrp argument. The leaves() method is a wrapper for the lgrp\_leaves() function.

#### isleaf(\$lgrp)

This method returns True if the Igroup specified by the value of the \$lgrp argument is a leaf Igroup. A leaf Igroup has no children. This method returns False in all other cases. The isleaf method is a wrapper for the lgrp\_isleaf function.

#### latency(\$from, \$to)

This method returns the latency value between a hardware resource in the lgroup specified by the \$from argument to a hardware resource in the lgroup specified by the \$to argument. The latency method uses the lgrp\_latency() function in version 1 of liblgrp. The latency

method uses the lgrp\_latency\_cookie() function in newer versions of liblgrp.



# Locality Group APIs

This chapter describes the APIs that applications use to interact with locality groups.

This chapter discusses the following topics:

- "Verifying the Interface Version" on page 31 describes the functions that give information about the interface.
- "Initializing the Locality Group Interface" on page 32 describes function calls that initialize and shut down the portion of the interface that is used to traverse the locality group hierarchy and to discover the contents of a locality group.
- "Locality Group Hierarchy" on page 33 describes function calls that navigate the locality group hierarchy and functions that get characteristics of the locality group hierarchy.
- "Locality Group Contents" on page 35 describes function calls that retrieve information about a locality group's contents.
- "Locality Group Characteristics" on page 37 describes function calls that retrieve information about a locality group's characteristics.
- "Locality Groups and Thread and Memory Placement" on page 38 describes how to affect the locality group placement of a thread and its memory.
- "Examples of API Usage" on page 44 contains code that performs example tasks by using the APIs that are described in this chapter.

### **Verifying the Interface Version**

The lgrp\_version(3LGRP) function must be used to verify the presence of a supported lgroup interface before using the lgroup API. The lgrp\_version() function has the following syntax:

```
#include <sys/lgrp_user.h>
int lgrp version(const int version);
```

The lgrp\_version() function takes a version number for the lgroup interface as an argument and returns the lgroup interface version that the system supports. When the current

implementation of the Igroup API supports the version number in the version argument, the Igrp\_version() function returns that version number. Otherwise, the Igrp\_version() function returns LGRP\_VER\_NONE.

### **Initializing the Locality Group Interface**

Applications must call <code>lgrp\_init(3LGRP)</code> in order to use the APIs for traversing the lgroup hierarchy and to discover the contents of the lgroup hierarchy. The call to <code>lgrp\_init()</code> gives the application a consistent snapshot of the lgroup hierarchy. The application developer can specify whether the snapshot contains only the resources that are available to the calling thread specifically or the resources that are available to the operating system in general. The <code>lgrp\_init()</code> function returns a cookie that is used for the following tasks:

- Navigating the lgroup hierarchy
- Determining the contents of an Igroup
- Determining whether the snapshot is current

### Using lgrp\_init()

The lgrp\_init() function initializes the lgroup interface and takes a snapshot of the lgroup hierarchy.

```
#include <sys/lgrp_user.h>
lgrp_cookie_t lgrp_init(lgrp_view_t view);
```

When the <code>lgrp\_init()</code> function is called with <code>LGRP\_VIEW\_CALLER</code> as the view, the function returns a snapshot that contains only the resources that are available to the calling thread. When the <code>lgrp\_init()</code> function is called with <code>LGRP\_VIEW\_OS</code> as the view, the function returns a snapshot that contains the resources that are available to the operating system. When a thread successfully calls the <code>lgrp\_init()</code> function, the function returns a cookie that is used by any function that interacts with the <code>lgroup</code> hierarchy. When a thread no longer needs the cookie, call the <code>lgrp\_fini()</code> function with the cookie as the argument.

The Igroup hierarchy consists of a root Igroup that contains all of the machine's CPU and memory resources. The root Igroup might contain other locality groups bounded by smaller latencies.

The lgrp\_init() function can return two errors. When a view is invalid, the function returns EINVAL. When there is insufficient memory to allocate the snapshot of the lgroup hierarchy, the function returns ENOMEM.

### Using lgrp\_fini()

The lgrp\_fini(3LGRP) function ends the usage of a given cookie and frees the corresponding lgroup hierarchy snapshot.

```
#include <sys/lgrp_user.h>
int lgrp_fini(lgrp_cookie_t cookie);
```

The lgrp\_fini() function takes a cookie that represents an lgroup hierarchy snapshot created by a previous call to lgrp\_init(). The lgrp\_fini() function frees the memory that is allocated to that snapshot. After the call to lgrp\_fini(), the cookie is invalid. Do not use that cookie again.

When the cookie passed to the lgrp fini() function is invalid, lgrp fini() returns EINVAL.

### **Locality Group Hierarchy**

The APIs that are described in this section enable the calling thread to navigate the lgroup hierarchy. The lgroup hierarchy is a directed acyclic graph that is similar to a tree, except that a node might have more than one parent. The root lgroup represents the whole machine and contains all of that machine's resources. The root lgroup is the lgroup with the highest latency value in the system. Each of the child lgroups contains a subset of the hardware that is in the root lgroup. Each child lgroup is bounded by a lower latency value. Locality groups that are closer to the root have more resources and a higher latency. Locality groups that are closer to the leaves have fewer resources and a lower latency. An lgroup can contain resources directly within its latency boundary. An lgroup can also contain leaf lgroups that contain their own sets of resources. The resources of leaf lgroups are available to the lgroup that encapsulates those leaf lgroups.

### Using lgrp\_cookie\_stale()

The lgrp\_cookie\_stale(3LGRP) function determines whether the snapshot of the lgroup hierarchy represented by the given cookie is current.

```
#include <sys/lgrp_user.h>
int lgrp_cookie_stale(lgrp_cookie_t cookie);
```

The cookie returned by the <code>lgrp\_init()</code> function can become stale due to several reasons that depend on the view that the snapshot represents. A cookie that is returned by calling the <code>lgrp\_init()</code> function with the view set to <code>LGRP\_VIEW\_OS</code> can become stale due to changes in

the lgroup hierarchy such as dynamic reconfiguration or a change in a CPU's online status. A cookie that is returned by calling the lgrp\_init() function with the view set to LGRP\_VIEW\_CALLER can become stale due to changes in the calling thread's processor set or changes in the lgroup hierarchy. A stale cookie is refreshed by calling the lgrp\_fini() function with the old cookie, followed by calling lgrp\_init() to generate a new cookie.

The lgrp cookie stale() function returns EINVAL when the given cookie is invalid.

### Using lgrp\_view()

The lgrp\_view(3LGRP) function determines the view with which a given lgroup hierarchy snapshot was taken.

```
#include <sys/lgrp_user.h>
lgrp_view_t lgrp_view(lgrp_cookie_t cookie);
```

The <code>lgrp\_view()</code> function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the snapshot's view of the lgroup hierarchy. Snapshots that are taken with the view <code>LGRP\_VIEW\_CALLER</code> contain only the resources that are available to the calling thread. Snapshots that are taken with the view <code>LGRP\_VIEW\_OS</code> contain all the resources that are available to the operating system.

The lgrp view() function returns EINVAL when the given cookie is invalid.

### Using lgrp\_nlgrps()

The lgrp\_nlgrps(3LGRP) function returns the number of locality groups in the system. If a system has only one locality group, memory placement optimizations have no effect.

```
#include <sys/lgrp_user.h>
int lgrp_nlgrps(lgrp_cookie_t cookie);
```

The lgrp\_nlgrps() function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the number of lgroups available in the hierarchy.

The lgrp\_nlgrps() function returns EINVAL when the cookie is invalid.

### Using lgrp\_root()

The lgrp root(3LGRP) function returns the root lgroup ID.

```
#include <sys/lgrp_user.h>
lgrp_id_t lgrp_root(lgrp_cookie_t cookie);
```

The lgrp\_root() function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the root lgroup ID.

### Using lgrp\_parents()

The lgrp\_parents(3LGRP) function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the number of parent lgroups for the specified lgroup.

If lgrp\_array is not NULL and the value of lgrp\_array\_size is not zero, the lgrp\_parents() function fills the array with parent lgroup IDs until the array is full or all parent lgroup IDs are in the array. The root lgroup has zero parents. When the lgrp\_parents() function is called for the root lgroup, lgrp\_array is not filled in.

The lgrp\_parents() function returns EINVAL when the cookie is invalid. The lgrp\_parents() function returns ESRCH when the specified lgroup ID is not found.

### Using lgrp\_children()

The lgrp\_children(3LGRP) function takes a cookie that represents the calling thread's snapshot of the lgroup hierarchy and returns the number of child lgroups for the specified lgroup.

If lgrp\_array is not NULL and the value of lgrp\_array\_size is not zero, the lgrp\_children() function fills the array with child lgroup IDs until the array is full or all child lgroup IDs are in the array.

The lgrp\_children() function returns EINVAL when the cookie is invalid. The lgrp\_children() function returns ESRCH when the specified lgroup ID is not found.

### **Locality Group Contents**

The following APIs retrieve information about the contents of a given lgroup.

The Igroup hierarchy organizes the domain's resources to simplify the process of locating the nearest resource. Leaf Igroups are defined with resources that have the least latency. Each of the successive ancestor Igroups of a given leaf Igroup contains the next nearest resources to its child Igroup. The root Igroup contains all of the resources that are in the domain.

The resources of a given lgroup are contained directly within that lgroup or indirectly within the leaf lgroups that the given lgroup encapsulates. Leaf lgroups directly contain their resources and do not encapsulate any other lgroups.

### Using lgrp\_resources()

The lgrp\_resources() function returns the number of resources contained in a specified lgroup.

The <code>lgrp\_resources()</code> function takes a cookie that represents a snapshot of the <code>lgroup</code> hierarchy. That cookie is obtained from the <code>lgrp\_init()</code> function. The <code>lgrp\_resources()</code> function returns the number of resources that are in the <code>lgroup</code> with the ID that is specified by the value of the <code>lgrp</code> argument. The <code>lgrp\_resources()</code> function represents the resources with a set of <code>lgroups</code> that directly contain CPU or memory resources. The <code>lgrp\_rsrc\_t</code> argument can have the following two values:

```
LGRP_RSRC_CPU The lgrp_resources() function returns the number of CPU resources.

LGRP_RSRC_MEM The lgrp_resources() function returns the number of memory resources.
```

When the value passed in the <code>lgrpids[]</code> argument is not null and the count argument is not zero, the <code>lgrp\_resources()</code> function stores <code>lgroup</code> IDs in the <code>lgrpids[]</code> array. The number of <code>lgroup</code> IDs stored in the array can be up to the value of the count argument.

The <code>lgrp\_resources()</code> function returns <code>EINVAL</code> when the specified cookie, <code>lgroup ID</code>, or type are not valid. The <code>lgrp\_resources()</code> function returns <code>ESRCH</code> when the function does not find the specified <code>lgroup ID</code>.

### Using lgrp cpus()

The lgrp\_cpus(3LGRP) function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the number of CPUs in a given lgroup.

If the *cpuid[]* argument is not NULL and the CPU count is not zero, the lgrp\_cpus() function fills the array with CPU IDs until the array is full or all the CPU IDs are in the array.

The *content* argument can have the following two values:

```
LGRP_CONTENT_ALL

The lgrp_cpus() function returns IDs for the CPUs in this lgroup and this lgroup's descendants.

LGRP_CONTENT_DIRECT

The lgrp_cpus() function returns IDs for the CPUs in this lgroup only.
```

The lgrp\_cpus() function returns EINVAL when the cookie, lgroup ID, or one of the flags is not valid. The lgrp\_cpus() function returns ESRCH when the specified lgroup ID is not found.

## Using lgrp\_mem\_size()

The lgrp\_mem\_size(3LGRP) function takes a cookie that represents a snapshot of the lgroup hierarchy and returns the size of installed or free memory in the given lgroup. The lgrp\_mem\_size() function reports memory sizes in bytes.

The *type* argument can have the following two values:

LGRP\_MEM\_SZ\_FREE The lgrp\_mem\_size() function returns the amount of free

memory in bytes.

LGRP\_MEM\_SZ\_INSTALLED The lgrp\_mem\_size() function returns the amount of installed

memory in bytes.

The *content* argument can have the following two values:

LGRP\_CONTENT\_ALL The lgrp\_mem\_size() function returns the amount of memory in

this Igroup and this Igroup's descendants.

LGRP CONTENT DIRECT The lgrp mem size() function returns the amount of memory in

this lgroup only.

The <code>lgrp\_mem\_size()</code> function returns <code>EINVAL</code> when the cookie, <code>lgroup ID</code>, or one of the flags is not valid. The <code>lgrp\_mem\_size()</code> function returns <code>ESRCH</code> when the specified <code>lgroup ID</code> is not found.

# **Locality Group Characteristics**

The following API retrieves information about the characteristics of a given lgroup.

## Using lgrp\_latency\_cookie()

The lgrp\_latency(3LGRP) function returns the latency between a CPU in one lgroup to the memory in another lgroup.

```
#include <sys/lgrp_user.h>
int lgrp_latency_cookie(lgrp_cookie_t cookie, lgrp_id_t from, lgrp_id_t to.
```

lat\_between\_t between);

The <code>lgrp\_latency\_cookie()</code> function takes a cookie that represents a snapshot of the <code>lgroup</code> hierarchy. The <code>lgrp\_init()</code> function creates this cookie. The <code>lgrp\_latency\_cookie()</code> function returns a value that represents the latency between a hardware resource in the <code>lgroup</code> given by the value of the <code>from</code> argument and a hardware resource in the <code>lgroup</code> given by the value of the <code>to</code> argument. If both arguments point to the same <code>lgroup</code>, the <code>lgrp\_latency\_cookie()</code> function returns the <code>latency</code> value within that <code>lgroup</code>.

**Note** – The latency value returned by the <code>lgrp\_latency\_cookie()</code> function is defined by the operating system and is platform-specific. This value does not necessarily represent the actual latency between hardware devices. Use this value only for comparison within one domain.

When the value of the *between* argument is LGRP\_LAT\_CPU\_TO\_MEM, the lgrp\_latency\_cookie() function measures the latency from a CPU resource to a memory resource.

The <code>lgrp\_latency\_cookie()</code> function returns <code>EINVAL</code> when the <code>lgroup</code> ID is not valid. When the <code>lgrp\_latency\_cookie()</code> function does not find the specified <code>lgroup</code> ID, the "from" <code>lgroup</code> does not contain any CPUs, or the "to" <code>lgroup</code> does not have any memory, the <code>lgrp\_latency\_cookie()</code> function returns <code>ESRCH</code>.

# **Locality Groups and Thread and Memory Placement**

This section discusses the APIs used to discover and affect thread and memory placement with respect to lgroups.

- The lgrp\_home(3LGRP) function is used to discover thread placement.
- The meminfo(2) system call is used to discover memory placement.
- The MADV\_ACCESS flags to the madvise(3C) function are used to affect memory allocation among lgroups.
- The lgrp\_affinity\_set(3LGRP) function can affect thread and memory placement by setting a thread's affinity for a given lgroup.
- The affinities of an Igroup may specify an order of preference for Igroups from which to allocate resources.
- The kernel needs information about the likely pattern of an application's memory use in order to allocate memory resources efficiently.
- The madvise() function and its shared object analogue madv.so.1 provide this information to the kernel.

 A running process can gather memory usage information about itself by using the meminfo() system call.

## Using lgrp\_home()

The lgrp\_home() function returns the home lgroup for the specified process or thread.

```
#include <sys/lgrp_user.h>
lgrp_id_t lgrp_home(idtype_t idtype, id_t id);
```

The <code>lgrp\_home()</code> function returns <code>EINVAL</code> when the ID type is not valid. The <code>lgrp\_home()</code> function returns <code>EPERM</code> when the effective user of the calling process is not the superuser and the real or effective user ID of the calling process does not match the real or effective user ID of one of the threads. The <code>lgrp\_home()</code> function returns <code>ESRCH</code> when the specified process or thread is not found.

## Using madvise()

The madvise() function advises the kernel that a region of user virtual memory in the range starting at the address specified in *addr* and with length equal to the value of the *len* parameter is expected to follow a particular pattern of use. The kernel uses this information to optimize the procedure for manipulating and maintaining the resources associated with the specified range. Use of the madvise() function can increase system performance when used by programs that have specific knowledge of their access patterns over memory.

```
#include <sys/types.h>
#include <sys/mman.h>
int madvise(caddr t addr, size t len, int advice);
```

The madvise() function provides the following flags to affect how a thread's memory is allocated among lgroups:

MADV ACCESS DEFAULT	This flag resets the ker	nel's expected access	s pattern for the specified
	Ö	1	1 1

range to the default.

MADV ACCESS LWP This flag advises the kernel that the next LWP to touch the specified

address range is the LWP that will access that range the most. The kernel allocates the memory and other resources for this range and

the LWP accordingly.

MADV ACCESS MANY This flag advises the kernel that many processes or LWPs will access

the specified address range randomly across the system. The kernel

allocates the memory and other resources for this range

accordingly.

The madvise() function can return the following values:

EAGAIN Some or all of the mappings in the specified address range, from *addr* to *addr+len*,

are locked for I/O.

EINVAL The value of the *addr* parameter is not a multiple of the page size as returned by

sysconf(3C), the length of the specified address range is less than or equal to zero,

or the advice is invalid.

EIO An I/O error occurs while reading from or writing to the file system.

ENOMEM Addresses in the specified address range are outside the valid range for the address

space of a process or the addresses in the specified address range specify one or

more pages that are not mapped.

ESTALE The NFS file handle is stale.

## Using meminfo()

The meminfo() function gives the calling process information about the virtual memory and physical memory that the system has allocated to that process.

The meminfo() function can return the following types of information:

MEMINFO VPHYSICAL The physical memory address corresponding to the given virtual

address

MEMINFO VLGRP The Igroup to which the physical page corresponding to the given

virtual address belongs

MEMINFO VPAGESIZE The size of the physical page corresponding to the given virtual

address

MEMINFO VREPLCNT The number of replicated physical pages that correspond to the

given virtual address

MEMINFO VREPL | n The *n*th physical replica of the given virtual address

MEMINFO VREPL LGRP|n The Igroup to which the *n*th physical replica of the given virtual

address belongs

MEMINFO PLGRP The lgroup to which the given physical address belongs

The meminfo() function takes the following parameters:

inaddr An array of input addresses.

addr count The number of addresses that are passed to meminfo().

*info\_req* An array that lists the types of information that are being requested.

info\_count The number of pieces of information that are requested for each address in the

inaddr array.

outdata An array where the meminfo() function places the results. The array's size is

equal to the product of the values of the *info\_req* and *addr\_count* parameters.

validity An array of size equal to the value of the addr\_count parameter. The validity

array contains bitwise result codes. The 0th bit of the result code evaluates the validity of the corresponding input address. Each successive bit in the result code evaluates the validity of the response to the members of the *info\_req* array

in turn.

The meminfo() function returns EFAULT when the area of memory to which the *outdata* or *validity* arrays point cannot be written to. The meminfo() function returns EFAULT when the area of memory to which the *info\_req* or *inaddr* arrays point cannot be read from. The meminfo() function returns EINVAL when the value of *info\_count* exceeds 31 or is less than 1. The meminfo() function returns EINVAL when the value of *addr\_count* is less than zero.

**EXAMPLE 3-2** Use of meminfo() to Print Out Physical Pages and Page Sizes Corresponding to a Set of Virtual Addresses

```
print info(void **addrvec, int how many)
        static const int info[] = {
                MEMINFO VPHYSICAL,
                MEMINFO VPAGESIZE);
        uint64 t * inaddr = alloca(sizeof(uint64_t) * how_many);
        uint64 t * outdata = alloca(sizeof(uint64 t) * how many * 2;
        uint t * validity = alloca(sizeof(uint t) * how many);
        int i;
        for (i = 0; i < how many; i++)
                inaddr[i] = (uint64 t *)addr[i];
        if (meminfo(inaddr, how many, info,
                    sizeof (info)/ sizeof(info[0]),
                    outdata, validity) < 0)
        for (i = 0; i < how many; i++) {
                if (validity[i] & 1 == 0)
                        printf("address 0x%llx not part of address
                                        space\n",
                                inaddr[i]);
                else if (validity[i] & 2 == 0)
```

**EXAMPLE 3-2** Use of meminfo() to Print Out Physical Pages and Page Sizes Corresponding to a Set of Virtual Addresses (Continued)

# **Locality Group Affinity**

The kernel assigns a thread to a locality group when the lightweight process (LWP) for that thread is created. That Igroup is called the thread's *home Igroup*. The kernel runs the thread on the CPUs in the thread's home Igroup and allocates memory from that Igroup whenever possible. If resources from the home Igroup are unavailable, the kernel allocates resources from other Igroups. When a thread has affinity for more than one Igroup, the operating system allocates resources from Igroups chosen in order of affinity strength. Lgroups can have one of three distinct affinity levels:

- 1. LGRP\_AFF\_STRONG indicates strong affinity. If this lgroup is the thread's home lgroup, the operating system avoids rehoming the thread to another lgroup if possible. Events such as dynamic reconfiguration, processor, offlining, processor binding, and processor set binding and manipulation might still result in thread rehoming.
- 2. LGRP\_AFF\_WEAK indicates weak affinity. If this lgroup is the thread's home lgroup, the operating system rehomes the thread if necessary for load balancing purposes.
- 3. LGRP\_AFF\_NONE indicates no affinity. If a thread has no affinity to any lgroup, the operating system assigns a home lgroup to the thread.

The operating system uses Igroup affinities as advice when allocating resources for a given thread. The advice is factored in with the other system constraints. Processor binding and processor sets do not change Igroup affinities, but might restrict the Igroups on which a thread can run.

### Using lgrp\_affinity\_get()

The lgrp\_affinity\_get(3LGRP) function returns the affinity that a LWP has for a given lgroup.

```
#include <sys/lgrp_user.h>
lgrp_affinity_t lgrp_affinity_get(idtype_t idtype, id_t id, lgrp_id_t lgrp);
```

The *idtype* and *id* arguments specify the LWP that the lgrp\_affinity\_get() function examines. If the value of *idtype* is P\_PID, the lgrp\_affinity\_get() function gets the lgroup affinity for one of the LWPs in the process whose process ID matches the value of the *id* argument. If the value of *idtype* is P\_LWPID, the lgrp\_affinity\_get() function gets the lgroup affinity for the LWP of the current process whose LWP ID matches the value of the *id* argument. If the value of *idtype* is P\_MYID, the lgrp\_affinity\_get() function gets the lgroup affinity for the current LWP.

The <code>lgrp\_affinity\_get()</code> function returns <code>EINVAL</code> when the given <code>lgroup</code> or <code>ID</code> type is not valid. The <code>lgrp\_affinity\_get()</code> function returns <code>EPERM</code> when the effective user of the calling process is not the superuser and the <code>ID</code> of the calling process does not match the real or effective user <code>ID</code> of one of the <code>LWPs</code>. The <code>lgrp\_affinity\_get()</code> function returns <code>ESRCH</code> when a given <code>lgroup</code> or <code>LWP</code> is not found.

### Using lgrp\_affinity\_set()

The lgrp\_affinity\_set(3LGRP) function sets the affinity that a LWP or set of LWPs have for a given lgroup.

The *idtype* and *id* arguments specify the LWP or set of LWPs the <code>lgrp\_affinity\_set()</code> function examines. If the value of *idtype* is P\_PID, the <code>lgrp\_affinity\_set()</code> function sets the lgroup affinity for all of the LWPs in the process whose process ID matches the value of the *id* argument to the affinity level specified in the *affinity* argument. If the value of *idtype* is P\_LWPID, the <code>lgrp\_affinity\_set()</code> function sets the lgroup affinity for the LWP of the current process whose LWP ID matches the value of the *id* argument to the affinity level specified in the *affinity* argument. If the value of *idtype* is P\_MYID, the <code>lgrp\_affinity\_set()</code> function sets the lgroup affinity for the current LWP or process to the affinity level specified in the *affinity* argument.

The lgrp\_affinity\_set() function returns EINVAL when the given lgroup, affinity, or ID type is not valid. The lgrp\_affinity\_set() function returns EPERM when the effective user of the calling process is not the superuser and the ID of the calling process does not match the real or effective user ID of one of the LWPs. The lgrp\_affinity\_set() function returns ESRCH when a given lgroup or LWP is not found.

# **Examples of API Usage**

This section contains code for example tasks that use the APIs that are described in this chapter.

#### EXAMPLE 3-3 Move Memory to a Thread

The following code sample moves the memory in the address range between *addr* and *addr+len* near the next thread to touch that range.

```
#include <stdio.h>
#include <sys/mman.h>
#include <sys/types.h>

/*
   * Move memory to thread
   */
void
mem_to_thread(caddr_t addr, size_t len)
{
   if (madvise(addr, len, MADV_ACCESS_LWP) < 0)
        perror("madvise");
}</pre>
```

#### **EXAMPLE 3-4** Move a Thread to Memory

This sample code uses the meminfo() function to determine the lgroup of the physical memory backing the virtual page at the given address. The sample code then sets a strong affinity for that lgroup in an attempt to move the current thread near that memory.

```
#include <stdio.h>
#include <sys/lgrp user.h>
#include <sys/mman.h>
#include <sys/types.h>
 * Move a thread to memory
int
thread to memory(caddr t va)
    uint64 t
                addr;
    ulong t
               count;
    lgrp id t home;
    uint64 t
                lgrp;
    uint t
                request;
    uint t
               valid;
    addr = (uint64 t)va;
    count = 1;
    request = MEMINFO VLGRP;
    if (meminfo(&addr, 1, &request, 1, &lgrp, &valid) != 0) {
```

### **EXAMPLE 3-4** Move a Thread to Memory (Continued)

```
perror("meminfo");
        return (1);
   }
   if (lgrp affinity set(P LWPID, P MYID, lgrp, LGRP AFF STRONG) != 0) {
        perror("lgrp_affinity_set");
        return (2);
   }
   home = lgrp home(P LWPID, P MYID);
   if (home == -1) {
        perror ("lgrp_home");
        return (3);
   }
   if (home != lgrp)
        return (-1);
    return (0);
}
```

### **EXAMPLE 3–5** Walk the lgroup Hierarchy

The following sample code walks through and prints out the Igroup hierarchy.

```
#include <stdio.h>
#include <stdlib.h>
#include <sys/lgrp user.h>
#include <sys/types.h>
* Walk and print lgroup hierarchy from given lgroup
 * through all its descendants
*/
int
lgrp_walk(lgrp_cookie_t cookie, lgrp_id_t lgrp, lgrp_content_t content)
    lgrp affinity t
                       aff;
                       *children;
    lgrp id t
                       *cpuids;
    processorid t
                       i;
    int
                       ncpus;
    int
                       nchildren;
    int
                       nparents;
    lgrp id t
                       *parents;
    lgrp mem size t
                       size;
     * Print given lgroup, caller's affinity for lgroup,
     * and desired content specified
    printf("LGROUP #%d:\n", lgrp);
```

### **EXAMPLE 3–5** Walk the Igroup Hierarchy (Continued)

```
aff = lgrp affinity get(P LWPID, P MYID, lgrp);
if (aff == -1)
    perror ("lgrp affinity get");
printf("\tAFFINITY: %d\n", aff);
printf("CONTENT %d:\n", content);
 * Get CPUs
 */
ncpus = lgrp cpus(cookie, lgrp, NULL, 0, content);
printf("\t%d CPUS: ", ncpus);
if (ncpus == -1) {
    perror("lgrp cpus");
    return (-1);
} else if (ncpus > 0) {
    cpuids = malloc(ncpus * sizeof (processorid t));
    ncpus = lgrp_cpus(cookie, lgrp, cpuids, ncpus, content);
            if (ncpus == -1) {
        free(cpuids);
                       perror("lgrp cpus");
        return (-1);
    for (i = 0; i < ncpus; i++)
        printf("%d ", cpuids[i]);
    free(cpuids);
printf("\n");
* Get memory size
 */
printf("\tMEMORY: ");
size = lgrp mem size(cookie, lgrp, LGRP MEM SZ INSTALLED, content);
if (size == -1) {
    perror("lgrp mem size");
    return (-1);
}
printf("installed bytes 0x%llx, ", size);
size = lgrp mem size(cookie, lgrp, LGRP MEM SZ FREE, content);
    if (size == -1) {
    perror("lgrp mem size");
    return (-1);
}
printf("free bytes 0x%llx\n", size);
* Get parents
nparents = lgrp_parents(cookie, lgrp, NULL, 0);
printf("\t%d PARENTS: ", nparents);
if (nparents == -1) {
    perror("lgrp parents");
    return (-1);
} else if (nparents > 0) {
```

### **EXAMPLE 3–5** Walk the Igroup Hierarchy (Continued)

```
parents = malloc(nparents * sizeof (lgrp id t));
        nparents = lgrp parents(cookie, lgrp, parents, nparents);
                    if (nparents == -1) {
             free(parents);
                         perror("lgrp parents");
             return (-1);
        for (i = 0; i < nparents; i++)
    printf("%d ", parents[i]);</pre>
        free(parents);
    printf("\n");
     * Get children
    nchildren = lgrp children(cookie, lgrp, NULL, 0);
    printf("\t%d CHILDREN: ", nchildren);
    if (nchildren == -1) {
        perror("lgrp_children");
        return (-1);
    } else if (nchildren > 0) {
        children = malloc(nchildren * sizeof (lgrp_id_t));
        nchildren = lgrp_children(cookie, lgrp, children, nchildren);
                    if (nchildren == -1) {
            free(children);
                         perror("lgrp_children");
             return (-1);
        printf("Children: ");
        for (i = 0; i < nchildren; i++)
            printf("%d ", children[i]);
        printf("\n");
        for (i = 0; i < nchildren; i++)
            lgrp walk(cookie, children[i], content);
        free(children);
    printf("\n");
    return (0);
}
EXAMPLE 3-6 Find the Closest Igroup With Available Memory Outside a Given Igroup
#include <stdio.h>
#include <stdlib.h>
#include <sys/lgrp user.h>
#include <sys/types.h>
#define
           INT MAX
                       2147483647
```

EXAMPLE 3-6 Find the Closest Igroup With Available Memory Outside a Given Igroup (Continued)

```
* Find next closest lgroup outside given one with available memory
*/
lgrp id t
lgrp next nearest(lgrp cookie t cookie, lgrp id t from)
    lgrp_id_t
                       closest;
    int
                       i;
    int
                       latency;
    int
                       lowest;
    int
                       nparents;
    lgrp id t
                       *parents;
    lgrp mem size t
                       size;
     * Get number of parents
   nparents = lgrp_parents(cookie, from, NULL, 0);
    if (nparents == -1) {
        perror("lgrp parents");
        return (LGRP NONE);
    }
    * No parents, so current lgroup is next nearest
    if (nparents == 0) {
        return (from);
    }
     * Get parents
     */
   parents = malloc(nparents * sizeof (lgrp id t));
    nparents = lgrp_parents(cookie, from, parents, nparents);
    if (nparents == -1) {
        perror("lgrp_parents");
        free(parents);
        return (LGRP NONE);
     * Find closest parent (ie. the one with lowest latency)
   closest = LGRP NONE;
    lowest = INT MAX;
    for (i = 0; \bar{i} < nparents; i++) {
        lgrp id t lgrp;
         * See whether parent has any free memory
         */
        size = lgrp_mem_size(cookie, parents[i], LGRP MEM SZ FREE,
            LGRP CONTENT ALL);
```

**EXAMPLE 3-6** Find the Closest Igroup With Available Memory Outside a Given Igroup (Continued)

```
if (size > 0)
            lgrp = parents[i];
        else {
            if (size == -1)
                perror("lgrp mem size");
             * Find nearest ancestor if parent doesn't
             * have any memory
            lgrp = lgrp next nearest(cookie, parents[i]);
            if (lgrp == LGRP NONE)
                continue;
        }
         * Get latency within parent lgroup
        latency = lgrp_latency_cookie(lgrp, lgrp);
        if (latency == -1) {
            perror("lgrp_latency_cookie");
            continue;
        }
        * Remember lgroup with lowest latency
        if (latency < lowest) {</pre>
            closest = lgrp;
            lowest = latency;
        }
   }
    free(parents);
    return (closest);
}
* Find lgroup with memory nearest home lgroup of current thread
*/
lgrp id t
lgrp nearest(lgrp cookie t cookie)
{
    lgrp id t
                 home;
    longlong_t
                  size;
    * Get home lgroup
   home = lgrp home(P LWPID, P MYID);
     * See whether home lgroup has any memory available in its hierarchy
     */
```

EXAMPLE 3-6 Find the Closest Igroup With Available Memory Outside a Given Igroup (Continued)

### **EXAMPLE 3-7** Find Nearest Igroup With Free Memory

This example code finds the nearest lgroup with free memory to a given thread's home lgroup.

```
lgrp id t
lgrp nearest(lgrp cookie t cookie)
        lgrp id t
                          home:
        longlong t
                          size;
         * Get home lgroup
        home = lgrp home();
         * See whether home laroup has any memory available in its hierarchy
        if (lgrp mem size(cookie, lgrp, LGRP MEM SZ FREE,
            LGRP CONTENT ALL, &size) == -1)
                perror("lgrp mem size");
         * It does, so return the home lgroup.
        if (size > 0)
                return (home);
         * Otherwise, find next nearest lgroup outside of the home.
        return (lgrp next nearest(cookie, home));
}
```